

02 June 2002 (Adapted)

1.

	A	B	C	D	E	F
A	0	20	30	32	12	15
B	20	0	10	25	32	16
C	30	10	0	15	35	19
D	32	25	15	0	20	34
E	12	32	35	20	0	16
F	15	16	19	34	16	0

b) $A_{12} E_{16} F_{16} B_{10} C_{15} D_{32} A$ 101

c) AEFBCDA is not possible on the network
the tour would be

AEFBCDEA

d) $B_{10} C_{15} D_{20} E_{12} A_{15} F_{16} B$ 88

2)

	B1	B2	B3	B4
A1	-4	-5	-2	4
A2	-1	1	-1	2
A3	0	5	-2	-4
A4	-1	3	-1	1

Row min

-5

(-1)

-4

(-1)

Column Minimax (-)

=

Row Maximin (-)

∴ Stable solution exists

Column max

0 5 **(-1)** 4

b) Saddle points A2/B3 and A4/B3

c) $V(B) = 1$

3) minimax

Stage	State	Action	Dest.	Value
1	D	DT	T	$\max(8, X) = 8^*$
	E	ET	T	$\max(10, X) = 10^*$
	F	FT	T	$\max(6, X) = 6^*$
2	A	AD	D	$\max(7, 8) = 8^*$
		AE	E	$\max(8, 10) = 10$
	B	BE	E	$\max(9, 10) = 10$
		BF	F	$\max(3, 6) = 6^*$
	C	CE	E	$\max(6, 10) = 10$
		CF	F	$\max(9, 6) = 9^*$
3	S	SA	A	$\max(9, 8) = 9$
		SB	B	$\max(7, 6) = 7^*$
		SC	C	$\max(6, 9) = 9$

SBFT (7)

4)	B1	B2	B3
A1	3	5	4
A2	1	4	2
A3	6	3	7

- Row 1 Dominates Row 2 \therefore Row 2 can be deleted as player A should never play 2
- Column 3 dominates column 1 \therefore Column 3 can be deleted as player B should never play 3.

	B1	B2
A1	3	5
A3	6	3

let A play 1 prob = p
 A play 2 prob = $1-p$

if B plays 1	$V(A) = 3p + 6(1-p) = 6 - 3p$	0 1
B plays 2	$V(A) = 5p + 3(1-p) = 3 + 2p$	6 3
		3 5

$$6 - 3p = 3 + 2p \Rightarrow 5p = 3 \Rightarrow p = \frac{3}{5}$$

optimal play-safe strategy for A

play 1 prob = $\frac{3}{5}$ play 3 prob = $\frac{2}{5}$ $V(A) = \frac{21}{5}$.
 play 2 prob = 0!

Invert matrix so pay-off is for B.

	A1	A3	
B1	-3	-6	B plays 1 prob = q
B2	-5	-3	B plays 2 prob = $1-q$

if A plays 1 $V(B) = -3q - 6(1-q) = 2q - 6$
 A plays 3 $V(B) = -6q - 3(1-q) = -3q - 3$

$$2q - 6 = -3q - 3 \Rightarrow 5q = 3 \Rightarrow q = \frac{3}{5}$$

\therefore B should play 1 prob = $\frac{3}{5}$
 B should play 2 prob = $\frac{2}{5}$
 B should play 3 prob = 0

$$V(B) = 2\left(\frac{3}{5}\right) - 6 = -\frac{24}{5}$$

5)

	J1	J2	J3	J4
M1	14	5	8	7
M2	2	12	6	5
M3	7	8	3	9
M4	2	4	6	10

Reduce Rows

9	0	3	2	-5
0	10	4	3	-2
4	5	0	6	-3
0	2	4	8	-2

Reduce Columns

9	0	3	0
0	10	4	1
4	5	0	4
0	2	4	6
X	X	X	-2

3 lines
 \therefore not optimal
 Smallest uncovered
 = 1

10	0	4	0
0	9	4	0
4	4	0	3
0	1	4	5

4 lines needed
 to cover zeros
 \therefore optimal

	0		0
0			0
		0	
0			

M1 - J2 (5)
 M2 - J4 (5)
 M3 - J3 (3)
 M4 - J1 (2)

15

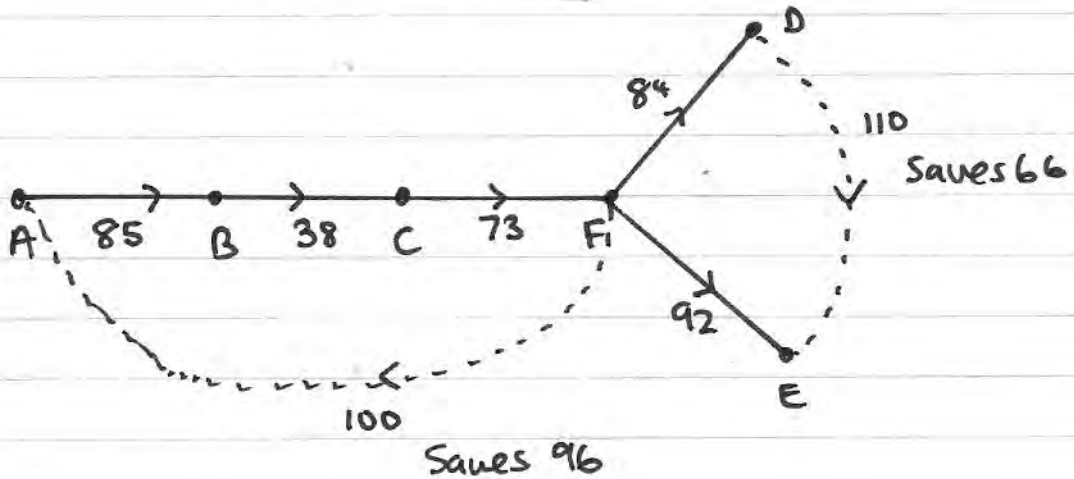
6)

	A	B	C	D	E	F
A	-	85	110	175	108	100
B	85	-	38	175	160	93
C	110	38	-	148	156	73
D	175	175	148	-	110	84
E	108	160	156	110	-	92
F	100	93	73	84	92	-

AB(85)
 BC(38)
 CF(73)
 FD(84)
 FE(92)

 372

\therefore Upper bound
 = 744



A₈₅ B₃₈ C₇₃ F₈₄ D₁₁₀ E₉₂ F₁₀₀ A 582

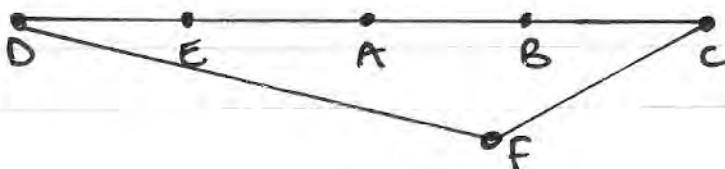
short cut so ABCFDEF CBA 678

using short cut DE.

c) Prim's

AB(85) RMST = 341
 BC(38) + FC(73)
 AE(108) + FD(84)
 ED(110) 498

Lower bound = 498



498 \leq optimal Solution \leq 582
 \therefore = 498

7) B_1 B_2 B_3 Supply

F_1	10	4	11	35
F_2	12	5	8	25
F_3	9	6	7	15

Demand 20 25 30 $\overline{175}$ balanced supply = demand

NWC

20	15		35
	10	15	25
		15	15
20	25	30	

Initial Solution

shadow costs

0	10	4	7
1	12	5	8
0	9	6	7

Improvement Indices

X	X	4
1	X	X
-1	2	X

$20 - \theta$	$15 + \theta$	
	$10 - \theta$	$15 + \theta$
$+\theta$		$15 - \theta$

entering Cell = $F_3 B_1$
 $\theta = 10$
 exiting Cell = $F_2 B_2$

10	25	
		25
10		5

	10	4	8
0	10	4	11
0	12	5	8
-1	9	6	7

Shadow Costs

X	X	3
2	1	X
X	3	X

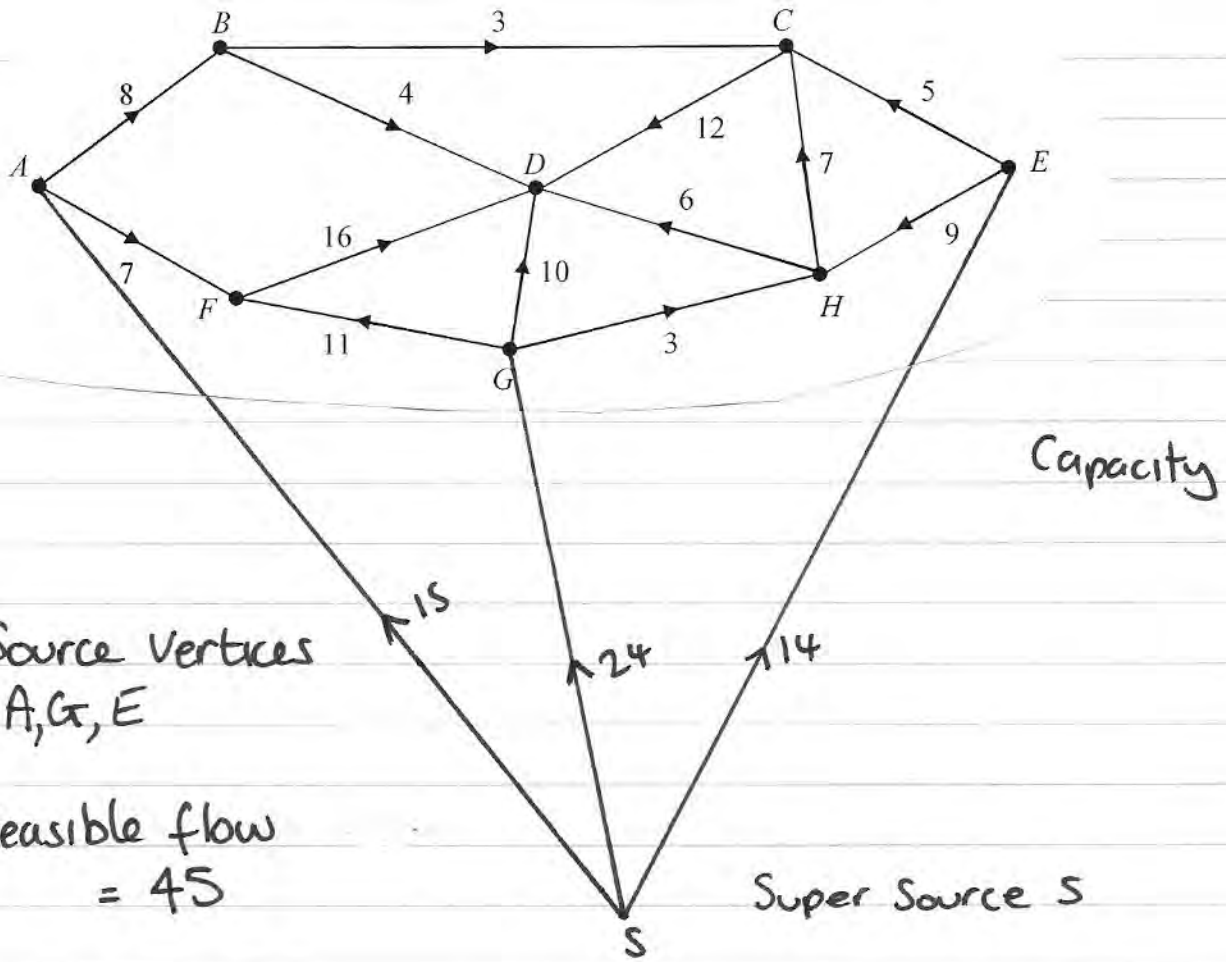
Improvement Indices

\therefore no negative Improvement Index \therefore Solution is optimal

	B_1	B_2	B_3
F_1	10	25	
F_2			25
F_3	10		5

Cost = 525

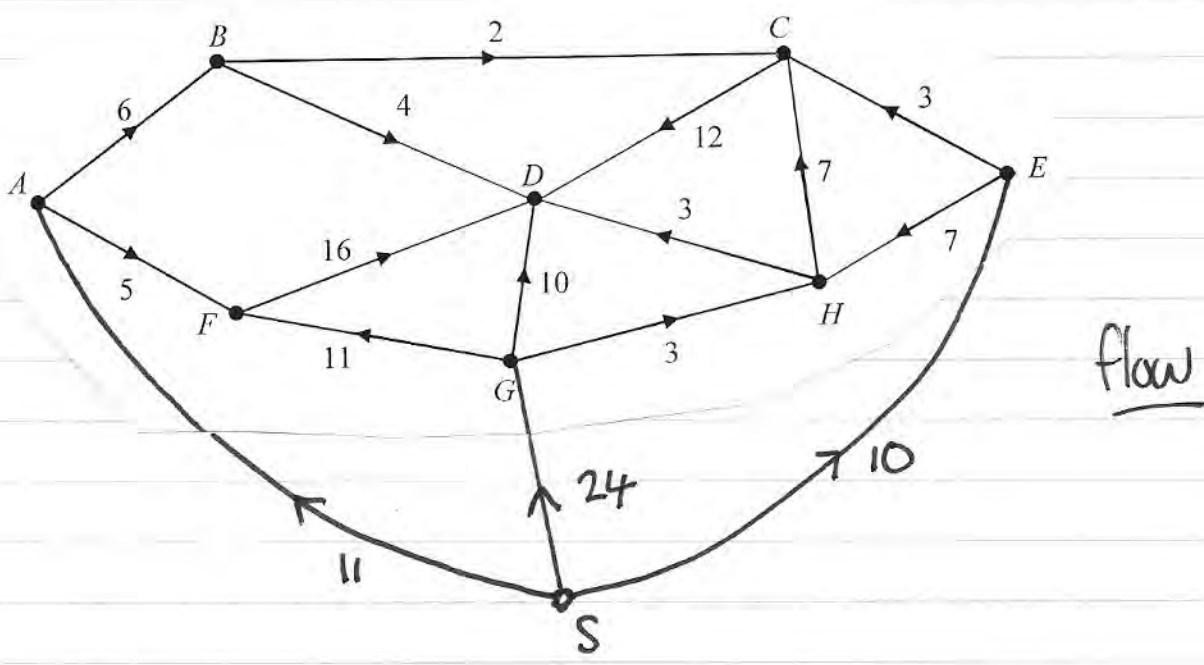
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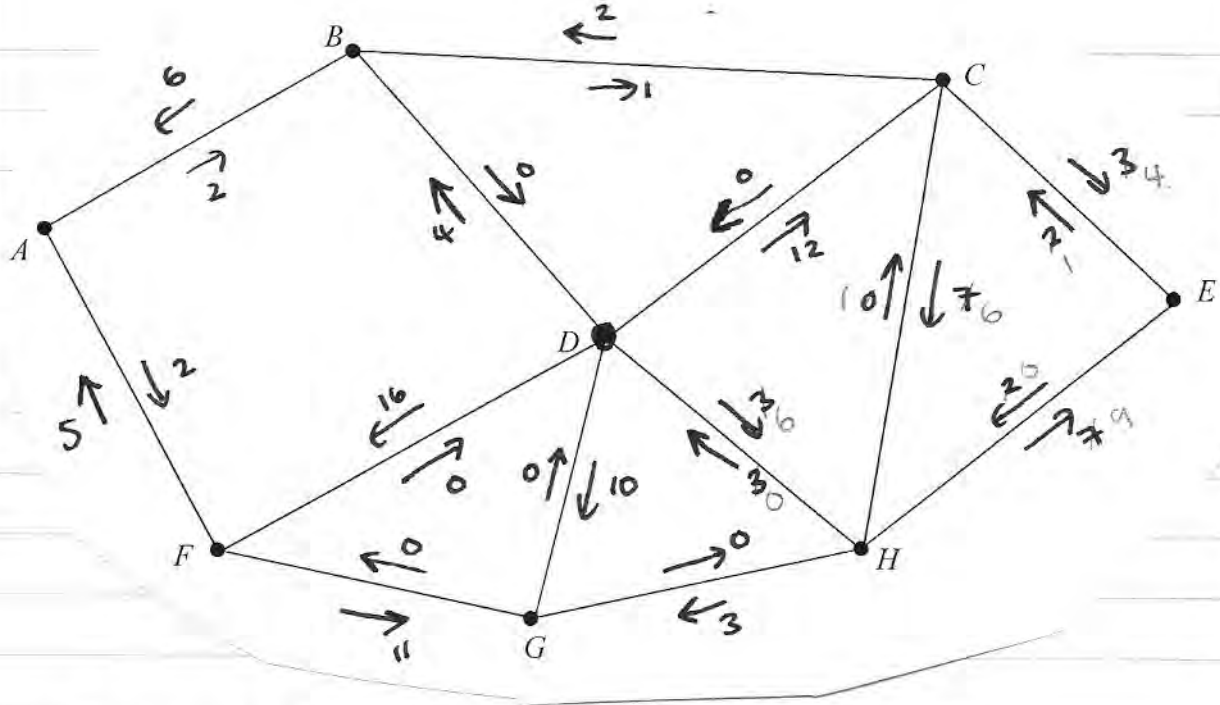


a) Source Vertices
A, G, E

b) feasible flow
= 45

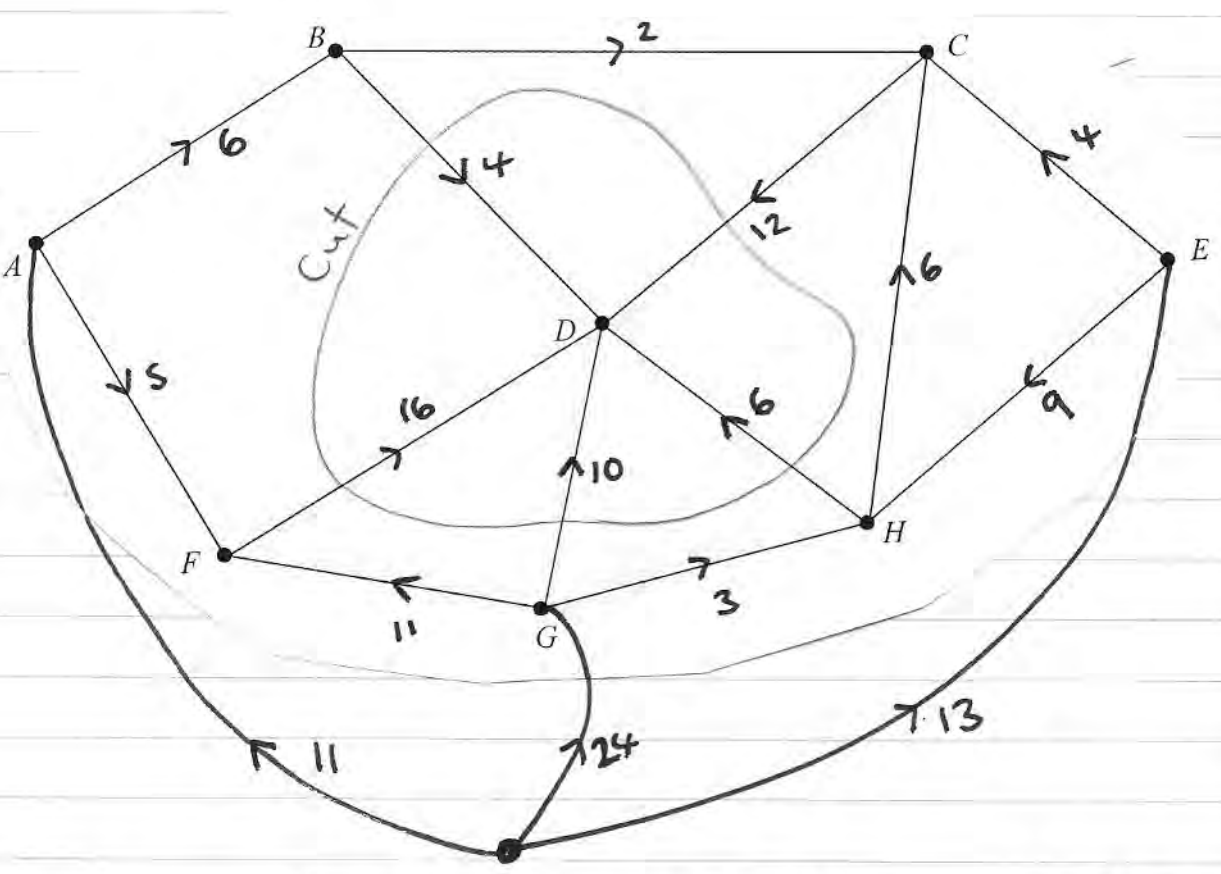
Super Source S





EHD(2)
ECHD(1)
3

total flow increased to 48



cut through BD, BC, DH, DG, DF splits network into 2 halves one containing sink the other the sources, all arcs saturated \therefore min cut = 48
 Max flow - min cut theorem \Rightarrow Max flow = 48 \therefore max

9)

	Processing	Blending	Packing	Profit (£100)	
x	Morning blend	3	1	2	4
y	Afternoon blend	2	3	4	5
z	Evening blend	4	2	3	3

≤ 35 ≤ 20 ≤ 24

Objective is to maximise profit P , (£100)

$P = 4x + 5y + 3z$

Subject to

$3x + 2y + 4z \leq 35$
 $x + 3y + 2z \leq 20$
 $2x + 4y + 3z \leq 24$
 $x, y, z \geq 0$

Basic variable	x	y	z	r	s	t	Value
r	3	2	4	1	0	0	35
s	1	3	2	0	1	0	20
t	2	4	3	0	0	1	24
P	-4	-5	-3	0	0	0	0

$0 = 35 \div 2 = 17.5$
 $0 = 20 \div 3 = 6\frac{2}{3}$
 $0 = 24 \div 4 = 6^*$

Increasing y

Basic variable	x	y	z	r	s	t	Value
	3	2	4	1	0	0	35
	1	3	2	0	1	0	20
	$\frac{1}{2}$	1	$\frac{3}{4}$	0	0	$\frac{1}{4}$	6
	-4	-5	-3	0	0	0	0

R_1
 R_2
 $R_3 \div 4$
 R_4

$Q = 11.5$
 $Q = -ve$
 $Q = 12$

Basic variable	x	y	z	r	s	t	Value	
r	2	0	5/2	1	0	-1/2	23	$R_1 - 2R_3$
S	-1/2	0	-1/4	0	1	-3/4	2	$R_2 - 3R_3$
y	1/2	1	3/4	0	0	1/4	6	R_3
P	-3/2	0	3/4	0	0	5/4	30	$R_4 + 5R_3$

Increasing x

Basic variable	x	y	z	r	s	t	Value	
	1	0	5/4	1/2	0	-1/4	23/2	$R_1 \div 2$
	-1/2	0	-1/4	0	1	-3/4	2	R_2
	1/2	1	3/4	0	0	1/4	6	R_3
	-3/2	0	3/4	0	0	5/4	30	R_4

Basic variable	x	y	z	r	s	t	Value	
x	1	0	5/4	1/2	0	-1/4	23/2	R_1
S	0	0	3/8	1/4	1	-7/8	31/4	$R_2 + \frac{1}{2}R_1$
y	0	1	1/8	-1/4	0	3/8	1/4	$R_3 - \frac{1}{2}R_1$
P	0	0	21/8	3/4	0	7/8	189/4	$R_4 + \frac{3}{2}R_1$

$$\begin{aligned}
 x &= 23/2 & r &= 0 & P &= 47.25 \\
 y &= 1/4 & s &= 31/4 & &= \pounds 4725 \\
 z &= 0 & t &= 0 & &
 \end{aligned}$$

c) Slack in S, which was used in the equation created from the constraints in the time for blending \therefore do not increase time available for blending but do increase time available for processing and packing as all the time is currently being used in this solution

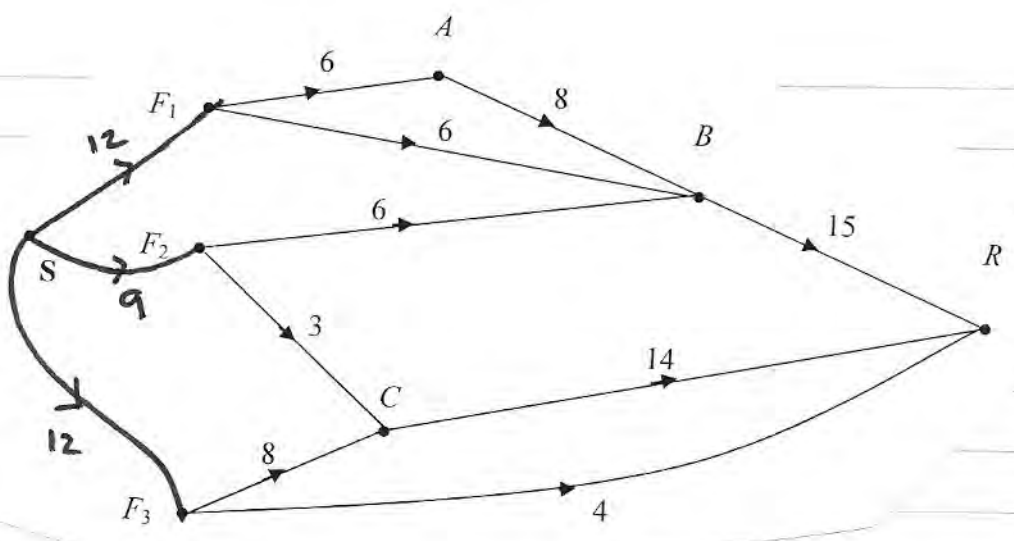
10) No negative values in the profit row \therefore optimal

$$\begin{array}{l} \text{b)} \quad x=1 \quad r=2/3 \quad P=11 \\ \quad \quad y=1/3 \quad s=0 \\ \quad \quad z=0 \quad t=0 \end{array}$$

$$\text{c)} \quad P+z+s+t=11 \Rightarrow P=11-z-s-t$$

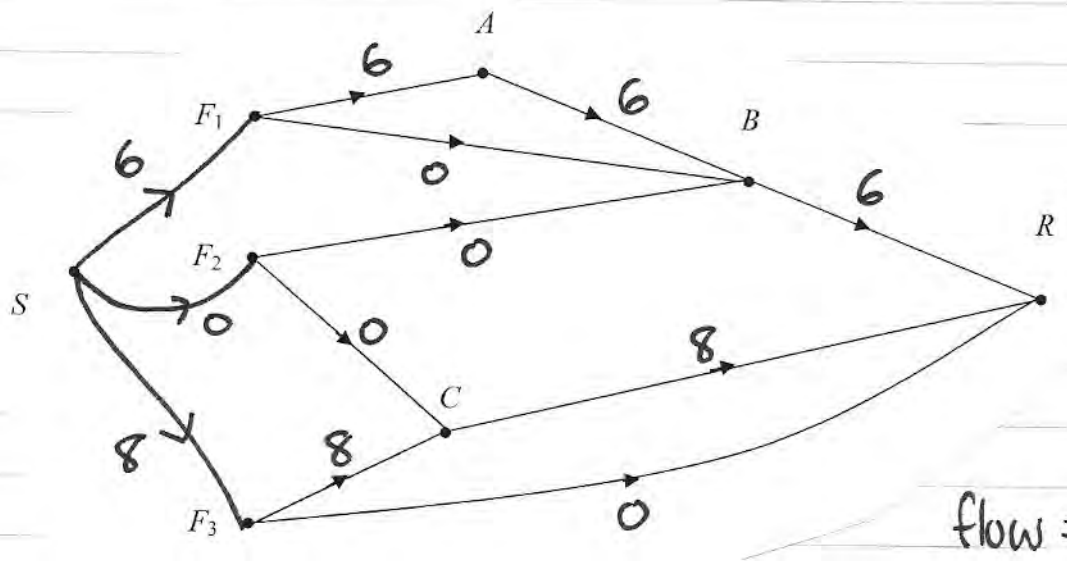
So increasing z, s or t would reduce the profit

11

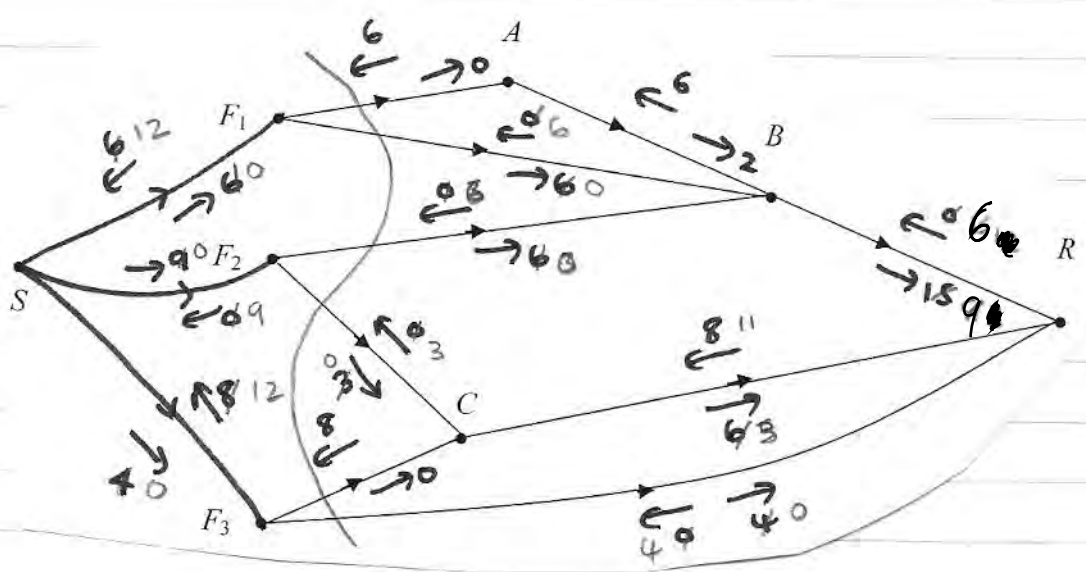


SF₁ABR = 6

SF₃CR = 8



flow = 14

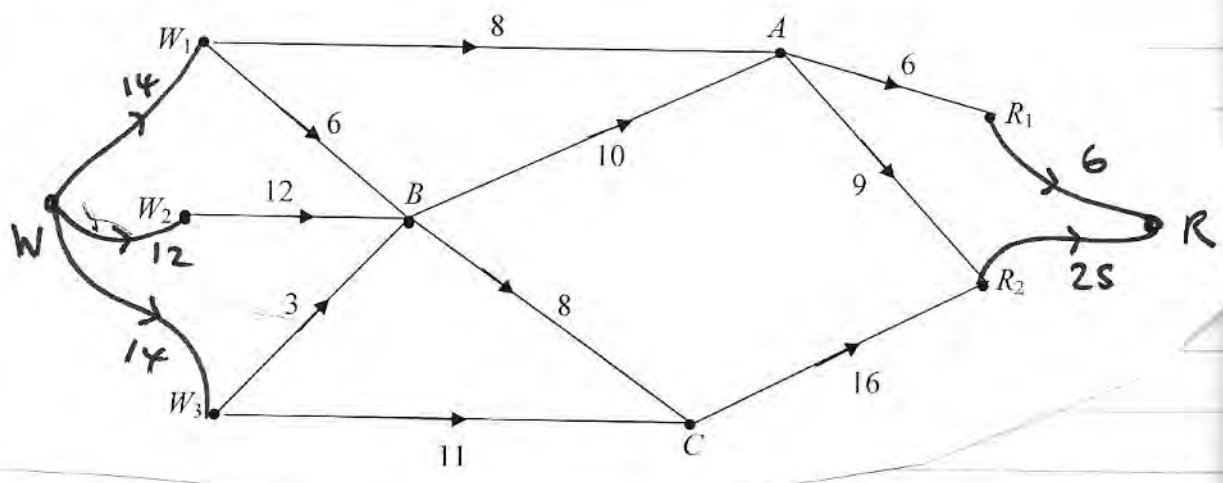


c) i) SF, BR (6)
SF₂BR (3)
SF₃R (4)
SF₂CR (3)
16

flow now 30

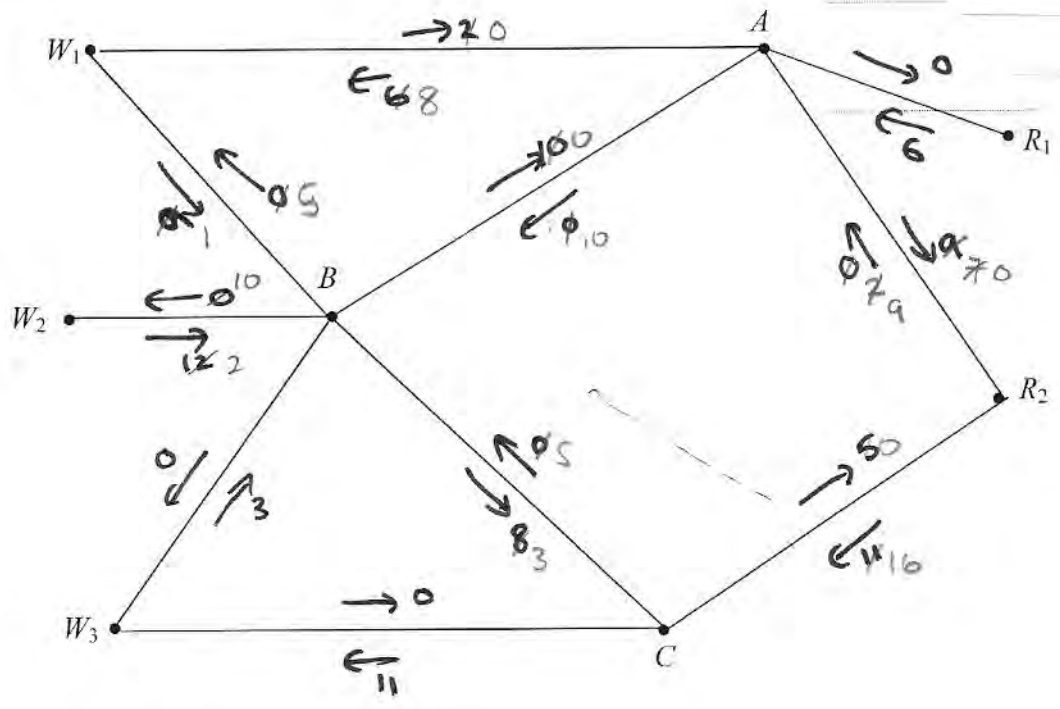
ii) Cut through $F_1A, F_1B, F_2B, F_2C, F_3C, F_3R$

only passes through saturated arcs \therefore min cut = 30 by maxflow-min cut theorem
max flow = 30 \therefore maximal.



$W W_1 A R_1 R = 6$

$W W_3 C R_2 R = 11$



- W W₁ A R₂ R (2)
- W W₁ B C R₂ R (5)
- W W₂ B A R₂ R (7)

flow now
31

d) flow through B = 15