

Exercise 4A

1 a

vertex	<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>	<i>F</i>
valency	2	3	2	3	3	3

There are 4 nodes with odd valency so the graph is *neither* Eulerian nor Semi-Eulerian.

b

vertex	<i>G</i>	<i>H</i>	<i>I</i>	<i>J</i>	<i>K</i>
valency	3	4	3	2	4

There are precisely 2 nodes of odd degree (*G* and *I*) so the graph is *semi-Eulerian*.
A possible route starting at *G* and finishing at *I* is: *GHKIJKGHI*

c

vertex	<i>L</i>	<i>M</i>	<i>N</i>	<i>P</i>	<i>Q</i>	<i>R</i>
valency	2	4	2	4	2	4

All vertices have even valency, so the graph is *Eulerian*.
A possible route starting and finishing at *L* is: *LMNPMRPQRL*

2 a i

vertex	<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>	<i>F</i>	<i>G</i>	<i>H</i>
valency	4	2	4	2	2	4	2	2

ii

vertex	<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>	<i>F</i>	<i>G</i>
valency	4	4	2	4	2	4	4

b i A possible route is: *ABCAFCEGHFDA*ii A possible route is: *ACFABEGBDGFDA*

3 a i

vertex	<i>R</i>	<i>S</i>	<i>T</i>	<i>U</i>	<i>V</i>	<i>W</i>
valency	2	2	3	3	2	2

Precisely 2 vertices of odd valency (*T* and *U*) so semi-Eulerian.

ii

vertex	<i>H</i>	<i>I</i>	<i>J</i>	<i>K</i>	<i>L</i>	<i>M</i>	<i>N</i>
valency	2	4	3	2	3	4	4

Precisely 2 nodes of odd degree (*J* and *L*) so semi-Eulerian.

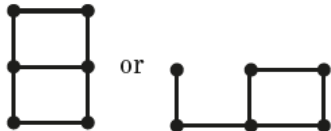
b i A possible route starting of *T* and finishing at *U* is: *TRSUWVTU*ii A possible route starts at *J* and finishes at *L*: *JKLMJIMNIHNL*

4 a The number of odd nodes of any graph must be even so this is not possible as there are 3 odd nodes.

4 b i $2x + 1 + 2x + 4x - 1 + 4x + 6x = 2E = 18$
 $\Rightarrow 18x = 18$
 $\Rightarrow x = 1$

ii Semi-Eulerian since there are two odd nodes.

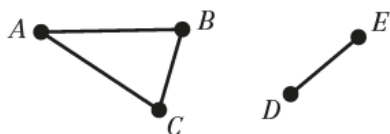
c Numerous possible answers e.g.:



5 a Not connected. There are no connections from A , B or C to D or E .

b Neither. To be Eulerian or semi-Eulerian the graph must be connected.

c



6 Adding up the numbers in each row, the orders of A , B , C , D , E are 2, 2, 2, 4, 4 respectively. Since they are all even the graph must be Eulerian.

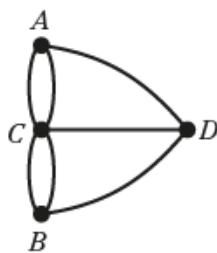
7 a n must be odd so that each vertex will have degree $n - 1$ which is even.



8 The example given in the question 1a is a counterexample. $ABEFCDA$ is a Hamiltonian cycle, but the graph is not Eulerian.

Challenge

a

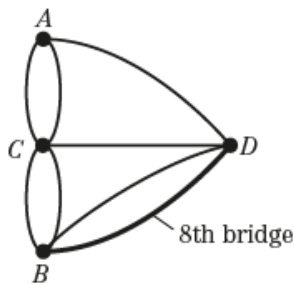


vertex	A	B	C	D
valency	3	3	5	3

There are more than two odd nodes, so the graph is *not* traversable.

Challenge

b

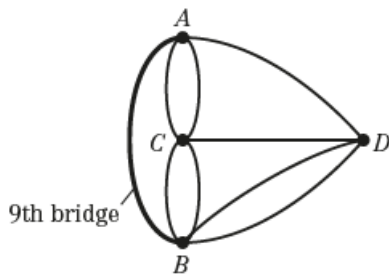


We will start at A and finish at C so these still need to have odd valency. We can only have two odd valencies so B and D must have even valencies (see table).

We need to change the valency of B and of D . So we build a bridge from B to D .

vertex	A	B	C	D
valency with 7 bridges	odd	odd	odd	odd
valency wanted	odd	even	odd	even

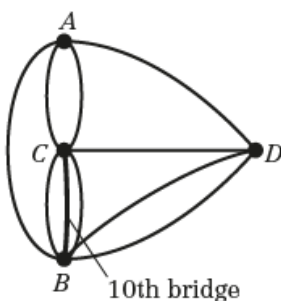
c



We will start at B and finish at C so these vertices need to be the two vertices with odd valency. We need A and D to have even valency (see table). We need to change the valency of node A and of node B . So we build a bridge from A to B .

vertex	A	B	C	D
valency with 8 bridges	odd	even	odd	even
valency wanted	even	odd	odd	even

d



All vertices now need to have even valency. This means we need to change the valencies of nodes B and C . So the 10th bridge needs to be built from B to C .

vertex	A	B	C	D
valency with 9 bridges	even	odd	odd	even
valency wanted	even	even	even	even